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Just-In-Time Compilation

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Definitions: Compiler

Definition: Compiler

Translation of source code from one programming language to another.

- ▶ Lexical Analysis
  Split the input into atomic units.

- ▶ Syntax Analysis
  Check if the syntax is correct (parser).

- ▶ Semantic Analysis
  E.g. type checking. Typically results in an intermediate representation.

- ▶ Optimization
  Increase the effectivity of the code without changing the semantics.

Definitions: AOT Compiler

Definition: Ahead-of-Time (AOT) Compiler

- Compilation before runtime
- Typically called by the programmer
- Often produces native code for direct execution
- Mostly platform dependent!

Definitions: Interpreters

Definition: Interpreter

- ► Does not create any output
- ► Processes at runtime
- ► Directly executes the code on the processor (=interpreting)
- ► Mostly platform independent!

Just-In-Time Compilation

Goals:

- ▶ Tries to fill the gap between Interpreter and AOT compilers
- ▶ Efficiency and platform independency

Structure:

- ▶ Performs compilation during runtime
- ▶ *Trace-based* or *method-based*
- ▶ Has to be as fast as possible
- ▶ Can use additional information for better optimizations
- ▶ Knows the underlying hardware

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Just-In-Time Compilation

Method-based:

- ► Only analyses the calls, no further analysis
- ► Similar to static AOT compilers
- ► Less efficient, but less overhead

Trace-based:

- ► Analyses what paths (*traces*) are often used, and to which methods they belong
- ► Can effectively be combined with an interpreter
- ► Allows targeted optimizations
- ► More efficient, more overhead

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Comparing JIT, AOT and Interpreters

AOT-Compilers:

+ Can be arbitrary slow and thus perform time-consuming optimizations
+ No overhead at runtime
− Little platform independence

Interpreters:

+ Platform independence
± Little overhead at runtime
− Inefficient for repetitive executions
− Little optimization possibilities

JIT-Compilers:

+ Platform independence
+ Highly optimized and efficient code
+ Runtime optimizations
− High overhead at runtime

Profiling

- ▶ Necessary for the optimization
- ▶ Analysis of the execution and collection of information
- ▶ A good profiling is crucial for efficient optimizations
- ▶ *static* or *dynamic*
- ▶ Has to produce as less overhead as possible (especially for JIT)

Next slides:

- ▶ How could a profiler be implemented?
- ▶ What information could be collected?

Profiling: How?

Sampling-based:

- ▶ Statistical approach: Take samples of the execution state
- ▶ Less precise, but very efficient
- ▶ Does not affect memory management much
- ▶ Drawback: A method could be completely undetected

Event-based:

- ▶ Profiler is triggered by certain events
- ▶ Very specified

Profiling: How?

Instrumentation:

- ▶ Injection of profiling code
- ▶ Very flexible, better performance than profiling by another thread

Profiling: What?

Call graph:

- ▶ Graph showing the call-dependencies of the methods
- ▶ Edges contain the amount of times this call was performed

Execution Trace:

- ▶ A trace representing the execution, including timestamps.
- ▶ Most precise profile

Optimization

Optimization:

- *Static* or *dynamic/adaptive*
- JIT compilers can use both types (but not every technique, because of their overhead)

Examples for static optimization:

- Dead code elimination
- Constant folding and propagation
- Loop-invariant code motion

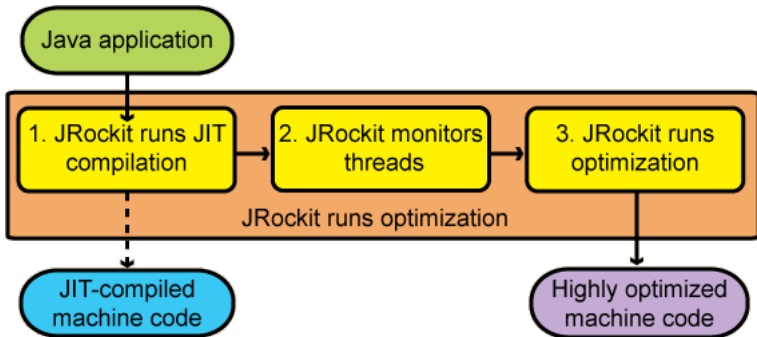Dynamic optimization on the basis of the *Java Virtual Machine*

JVM: Introduction

- ► Here: *HotSpot* and *JRockit*
- ► AOT compilation of Java source-code to Java bytecode
- ► JRockit uses JIT and HotSpot uses JIT& Interpreter
- ► Bytecode is assembly-like, and easy to process during runtime

JVM: Java Bytecode

What Java bytecode looks like

```
1  Compiled from "test.java"
2  public class test {
3    public test ();
4      Code:
5        0: aload_0
6        1: invokespecial #1                    // Method java/lang/Object."<
         init >":()V
7        4: return
8
9    public static void main(java.lang.String[]);
10     Code:
11       0: iconst_5
12       1: istore_1
13       2: return
14 }
```

JRockit structure

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JVM: HotSpot Optimizations

Hotspot Detection:

► Selective optimization: Only optimize parts of the code

► Assumption: the execution mostly resides in a small part of the code (80/20 rule)

► For 80% of the code , interpreting is probably more efficient than compiling

► Detect hot spots and focus on these for compilation and optimization

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JVM: HotSpot Optimizations

Method Inlining:

- ► Method invokations are time consuming
- ► Copy the code of frequently invokated methods inside their caller-methods
- ► Reduces overhead, allows more optimization

Problems:

- ► Recursive calls (infinite inlining, the optimizer has to set a maximum for the inlining of recursive calls)
- ► Overriding (see next slide)

JVM: HotSpot Optimizations

Dynamic deoptimization:

- ▶ Not every method can be inlined
- ▶ If a method is not final, it could possibly be overridden at runtime
- ▶ HotSpot speculatively inline methods which are not final (but not overridden at this moment)
- ▶ In certain cases, the inlining has to be undone, this is called dynamic deoptimization

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JVM: Hotspot Optimizations

Range check elimination:

- ▶ Java requires strict array bounding checks
- ▶ Reading, changing and overwriting a value would need two checks
- ▶ the JVM can check if it is possible for the index value to change between operations
- ▶ Reduces the range check amount

JVM: When to use JIT?

- As stated in HotSpot Detection, interpreting is in some cases better than JIT
- But: How to decide whether a method should be JIT-compiled or interpreted?
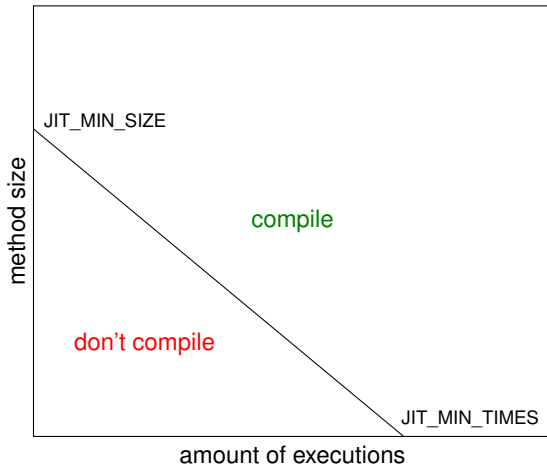- HotSpot uses heuristic approaches

JVM: When to use JIT?

The simple heuristics

- ► Often executed methods should be compiled
- ► Large methods should also be compiled, even when they are rarely executed
- ► Decision is based on the amount of executions and size of the methods

Example

Example for a simple decision graph:

Conclusion

- ▶ JIT compilation gives compilers the ability to be highly platform independent
- ▶ Optimizations mostly based on assumptions, which could also have negative effects
- ▶ Important decisions: Optimize? JIT or interpret?

Thank you for your attention.
Questions?